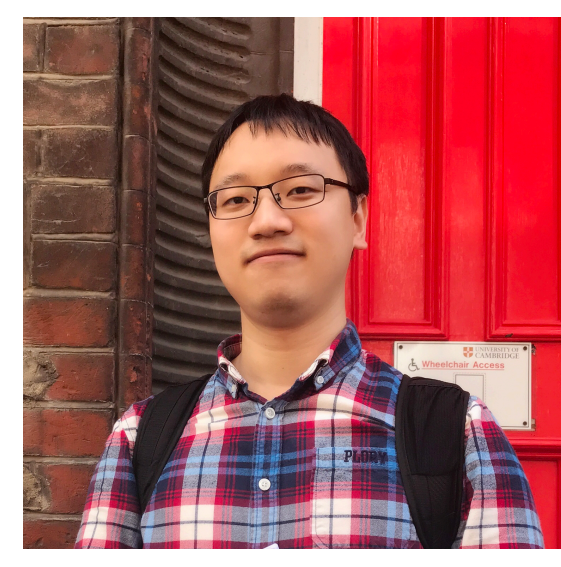
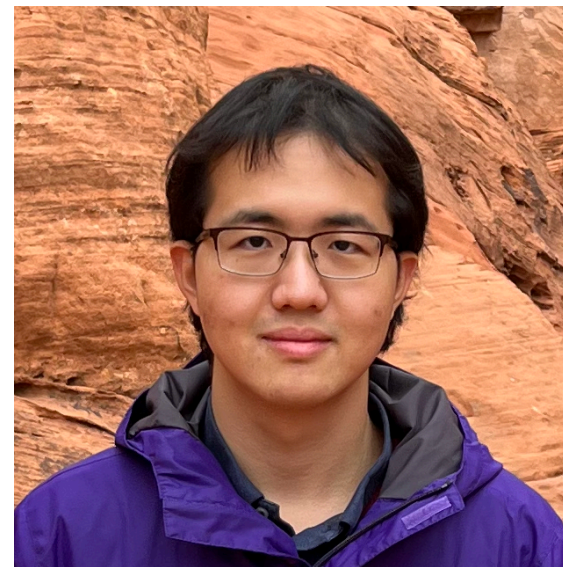


Draft, Sketch, and Prove

Sean Welleck

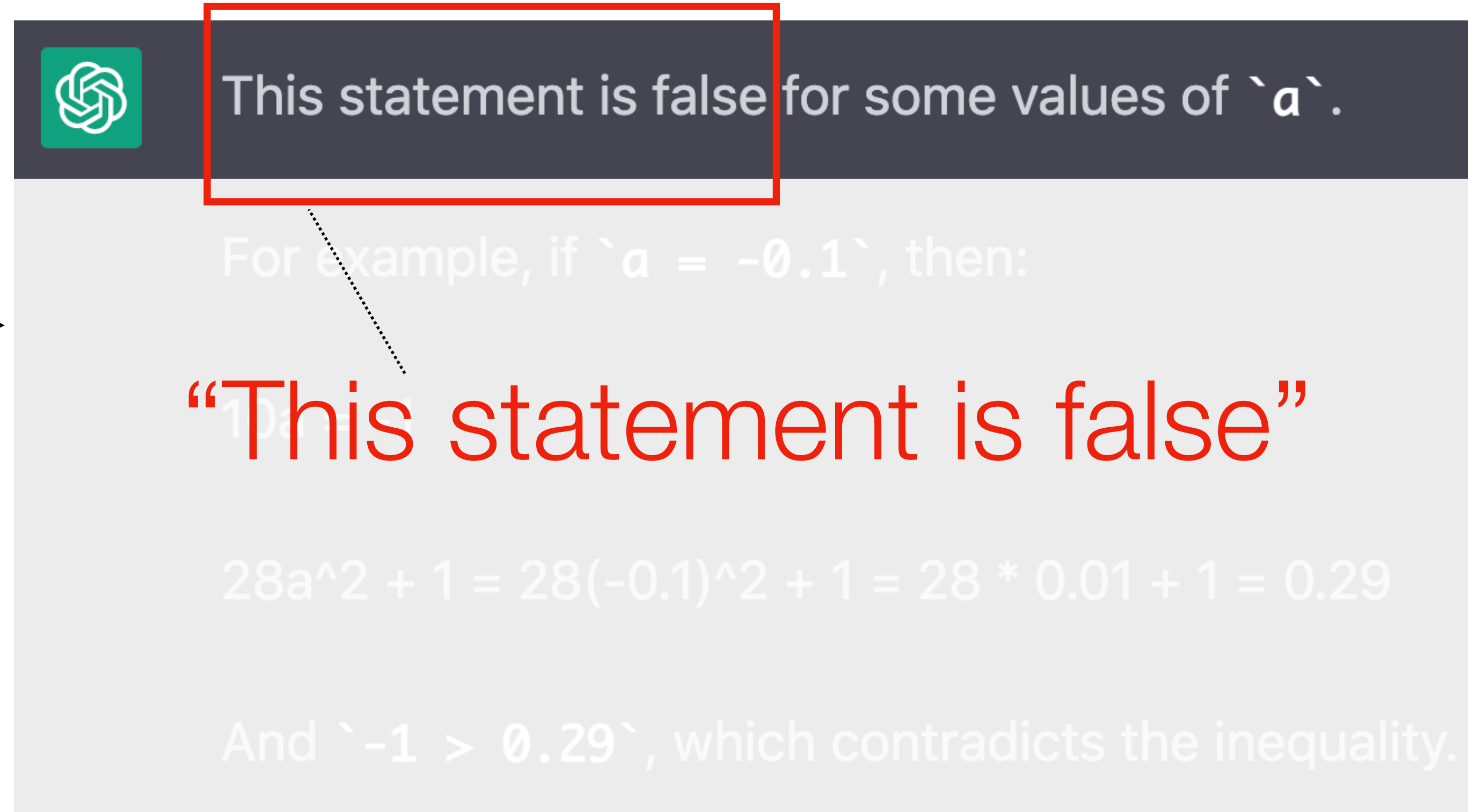
Draft, Sketch, and Prove: Guiding Formal Theorem Provers with Informal Proofs.

Jiang^{*}, Welleck^{*}, Zhou^{*}, Lacroix, Liu, Li, Jamnik, Lample⁺, Wu⁺
ICLR 2023



Informal mathematics

Show that for any
real number a ,
 $10a \leq 28a^2 + 1$

A screenshot of a Chat-GPT interface. At the top, a dark grey header bar contains the OpenAI logo and the text "This statement is false for some values of `a`." This header bar is highlighted with a red rectangular border. Below the header, the main content area is light grey. It contains the text "For example, if `a = -0.1`, then:" followed by the inequality $10a \leq 28a^2 + 1$ in red. Below this, the calculation $28a^2 + 1 = 28(-0.1)^2 + 1 = 28 * 0.01 + 1 = 0.29$ is shown. Finally, it states "And ` $-1 > 0.29$ `, which contradicts the inequality."

This statement is false for some values of `a`.

For example, if `a = -0.1`, then:

$10a \leq 28a^2 + 1$

$28a^2 + 1 = 28(-0.1)^2 + 1 = 28 * 0.01 + 1 = 0.29$

And ` $-1 > 0.29$ `, which contradicts the inequality.

Informal mathematics

Show that for any
real number a ,
 $10a \leq 28a^2 + 1$



This statement is false for some values of a .

For example, if $a = -0.1$, then:

$$10a = -1$$

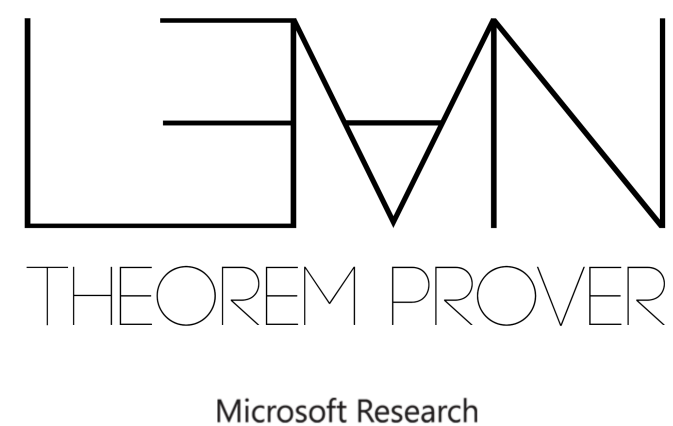
$$28a^2 + 1 = 28(-0.1)^2 + 1 = 28 * 0.01 + 1 = 0.29$$

“And $-1 > 0.29$ ”

And $-1 > 0.29$, which contradicts the inequality.

Formalized mathematics

- Translate mathematics into “code”, grounded in logic
 - Verified correctness
 - New ways of collaborating, teaching, thinking



Declarative proof

```
have c1: "10*280 = n*40"  
using assms
```

```
  by (smt (z3) prod_gcd_lcm_nat)
```

```
then have c2: "n = 10*280/40"
```

```
  by auto
```

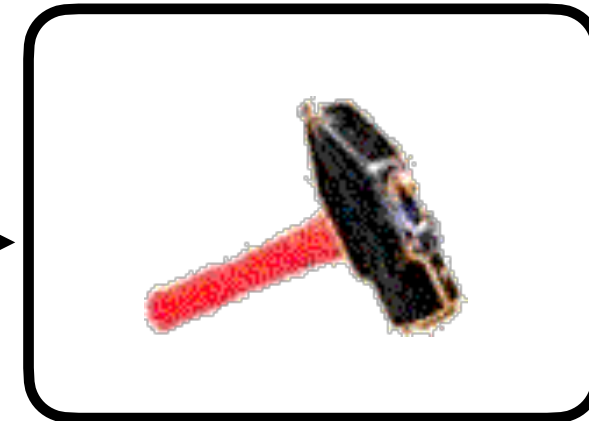
```
then show ?thesis
```

```
  by auto
```



Proof automation

```
have c1: "10*280 = n*40"  
using assms
```



```
have c1: "10*280 = n*40"  
using assms
```

```
by (smt (z3) prod_gcd_lcm_nat)
```

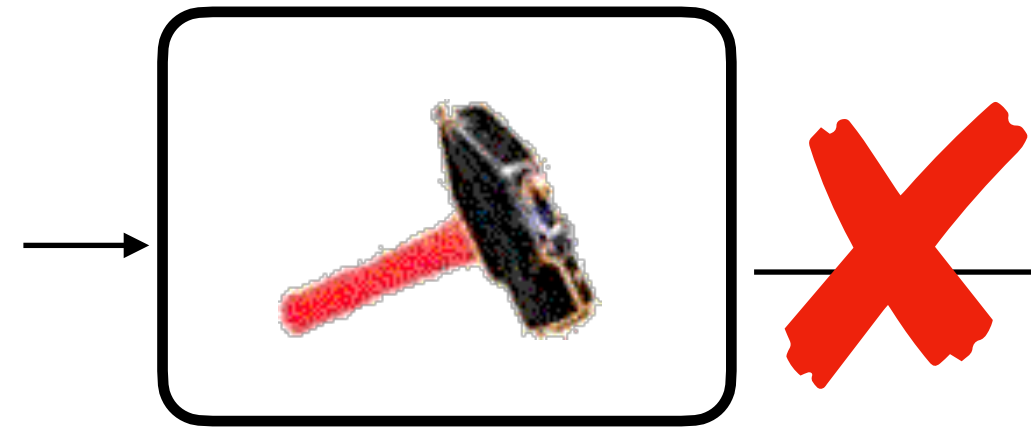
Sledgehammer [Paulson 2010]

- calls out to external provers

Proof automation

```
theorem algebra_absx1pabsxpabsxp1eqxp2_0leqxleq1:  
  fixes x :: real  
  assumes "abs (x - 1) + abs x + abs (x + 1) = x + 2"  
  shows "0 <= x <= 1"
```

Theorem



```
theorem algebra_absx1pabsxpabsxp1eqxp2_0leqxleq1:  
  fixes x :: real  
  assumes "abs (x - 1) + abs x + abs (x + 1) = x + 2"  
  shows "0 <= x <= 1"  
proof -  
  have c0: "x <= -1 <=> False"  
  proof -  
    assume c1: "x <= -1"  
    have c2: "abs(x-1) + abs x + abs(x+1) = -(x-1) - x - (x+1)" using c1  
    then have c3: "abs(x-1) + abs x + abs(x+1) = -3*x" by auto </ATP>  
    then have c4: "-3*x = x+2" using assms c3 <ATP> by auto </ATP>  
    then have c5: "x = -1/2" <ATP> by auto </ATP>  
    then show ?thesis using c1 <ATP> by auto </ATP>  
  qed  
  have c6: "-1 < x <=> x < 0 <=> False"  
  proof -  
    assume c7: "-1 < x" assume c8: "x < 0"  
    have c9: "abs(x-1) + abs x + abs(x+1) = -(x-1) - x + (x+1)" using c7 c8  
    <ATP> by auto </ATP>  
    then have c10: "abs(x-1) + abs x + abs(x+1) = 2-x" <ATP> by auto </ATP>  
    then have c11: "2-x = x+2" using assms c10 <ATP> by auto </ATP>  
    then have c12: "x = 0" <ATP> by auto </ATP>  
    then show ?thesis using c8 <ATP> by auto </ATP>  
  qed  
  have c13: "x > 1 <=> False"  
  proof -  
    assume c14: "x > 1"  
    have c15: "abs(x-1) + abs x + abs(x+1) = x-1 + x + (x+1)" using c14  
    <ATP> by auto </ATP>  
    then have c16: "abs(x-1) + abs x + abs(x+1) = 3*x" <ATP> by auto </ATP>  
    then have c17: "3*x = x+2" using assms c16 <ATP> by auto </ATP>  
    then have c18: "x = 1" <ATP> by auto </ATP>  
    then show ?thesis using c14 <ATP> by auto </ATP>  
  qed  
  then show ?thesis using c0 c6 c13 <ATP> by fastforce </ATP>  
qed
```

complex proof

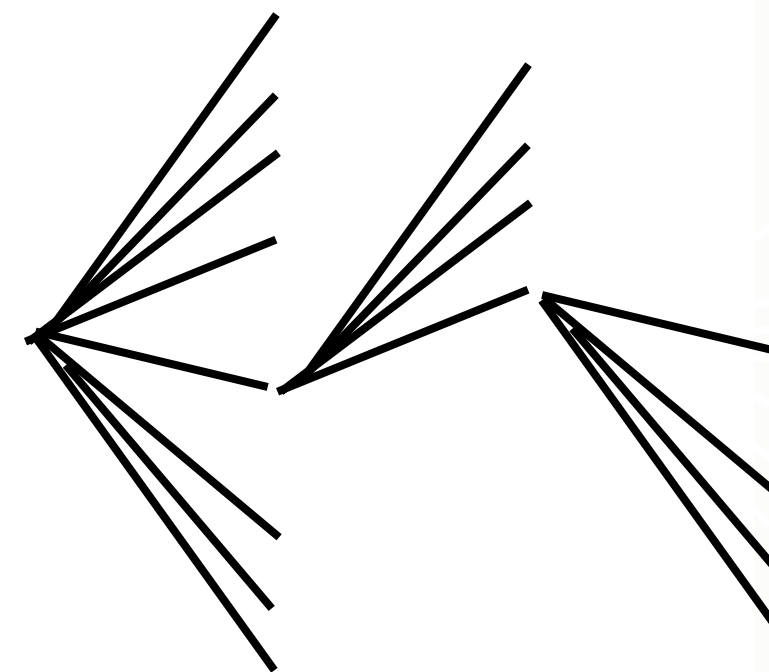
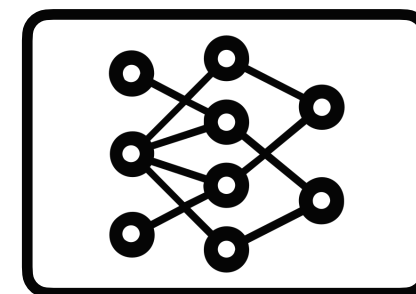
Sledgehammer [Paulson 2010]

Neural proof automation

- Train neural network on (context, next-step) pairs
- Tree search using next-step suggestions

```
theorem algebra_absx1pabsxpabsxp2_0leqxleq1:  
fixes x ::real assumes "abs (x - 1) + abs x + abs (x + 1) = x + 2"  
shows "0 <= x <= 1"
```

Theorem



```
theorem algebra_absx1pabsxpabsxp2_0leqxleq1:  
fixes x ::real assumes "abs (x - 1) + abs x + abs (x + 1) = x + 2"  
shows "0 <= x <= 1"  
proof -  
  have c0: "x <= -1 <math>\Longrightarrow</math> False"  
  proof -  
    assume c1: "x <= -1"  
    have c2: "abs(x-1) + abs x + abs(x+1) = -(x-1) - x - (x+1)" using c1  
      <math>\langle ATP \rangle</math> by auto </math></math>  
    then have c3: "abs(x-1) + abs x + abs(x+1) = -3*x" <math>\langle ATP \rangle</math> by auto </math></math>  
    then have c4: "-3*x = x+2" using assms c3 <math>\langle ATP \rangle</math> by auto </math></math>  
    then have c5: "x = -1/2" <math>\langle ATP \rangle</math> by auto </math></math>  
    then show ?thesis using c1 <math>\langle ATP \rangle</math> by auto </math></math>  
  qed  
  have c6: "-1 < x <math>\Longrightarrow</math> x < 0 <math>\Longrightarrow</math> False"  
  proof -  
    assume c7: "-1 < x" assume c8: "x < 0"  
    have c9: "abs(x-1) + abs x + abs(x+1) = -(x-1) - x + (x+1)" using c7 c8  
      <math>\langle ATP \rangle</math> by auto </math></math>  
    then have c10: "abs(x-1) + abs x + abs(x+1) = 2-x" <math>\langle ATP \rangle</math> by auto </math></math>  
    then have c11: "2-x = x+2" using assms c10 <math>\langle ATP \rangle</math> by auto </math></math>  
    then have c12: "x = 0" <math>\langle ATP \rangle</math> by auto </math></math>  
    then show ?thesis using c8 <math>\langle ATP \rangle</math> by auto </math></math>  
  qed  
  have c13: "x > 1 <math>\Longrightarrow</math> False"  
  proof -  
    assume c14: "x > 1"  
    have c15: "abs(x-1) + abs x + abs(x+1) = x-1 + x + (x+1)" using c14  
      <math>\langle ATP \rangle</math> by auto </math></math>  
    then have c16: "abs(x-1) + abs x + abs(x+1) = 3*x" <math>\langle ATP \rangle</math> by auto </math></math>  
    then have c17: "3*x = x+2" using assms c16 <math>\langle ATP \rangle</math> by auto </math></math>  
    then have c18: "x = 1" <math>\langle ATP \rangle</math> by auto </math></math>  
    then show ?thesis using c14 <math>\langle ATP \rangle</math> by auto </math></math>  
  qed  
  then show ?thesis using c0 c6 c13 <math>\langle ATP \rangle</math> by fastforce </math></math>  
qed
```

Neural proof automation

- Train neural network on (context, next-step) pairs
- Tree search using next-step candidates
- Limited amount of formal data
- Large search space -> expensive
 - Smaller network
 - Smaller # of candidates

Informal \rightarrow formal proof automation

Informal
Proof



Formal
Proof

```
theorem algebra_absxmpabsxpabsxp2_01eqxleq1:
  fixes x ::real assumes "abs (x - 1) + abs x + abs (x + 1) = x + 2"
  shows "0 <|e> x <and> x <|e> 1"
proof -
  have c0: "x <|e> -1 <Longrightarrow> False"
  proof -
    assume c1: "x <|e> -1"
    have c2: "abs(x-1) + abs x + abs(x+1) = -(x-1) - x - (x+1)" using c1
      <ATP> by auto </ATP>
    then have c3: "abs(x-1) + abs x + abs(x+1) = -3*x" <ATP> by auto </ATP>
    then have c4: "-3*x = x+2" using assms c3 <ATP> by auto </ATP>
    then have c5: "x = -1/2" <ATP> by auto </ATP>
    then show ?thesis using c1 <ATP> by auto </ATP>
  qed
  have c6: "-1 <|e> x <Longrightarrow> False"
  proof -
    assume c7: "-1 < x"
    assume c8: "x < 0"
    have c9: "abs(x-1) + abs x + abs(x+1) = -(x-1) - x + (x+1)" using c7 c8
      <ATP> by auto </ATP>
    then have c10: "abs(x-1) + abs x + abs(x+1) = 2-x" <ATP> by auto </ATP>
    then have c11: "2-x = x+2" using assms c10 <ATP> by auto </ATP>
    then have c12: "x = 0" <ATP> by auto </ATP>
    then show ?thesis using c7 <ATP> by auto </ATP>
  qed
  have c13: "x > 1 <Longrightarrow> False"
  proof -
    assume c14: "x > 1"
    have c15: "abs(x-1) + abs x + abs(x+1) = x-1 + x + (x+1)" using c14
      <ATP> by auto </ATP>
    then have c16: "abs(x-1) + abs x + abs(x+1) = 3*x" <ATP> by auto </ATP>
    then have c17: "3*x = x+2" using assms c16 <ATP> by auto </ATP>
    then have c18: "x = 1" <ATP> by auto </ATP>
    then show ?thesis using c14 <ATP> by auto </ATP>
  qed
  then show ?thesis using c0 c6 c13 <ATP> by fastforce </ATP>
qed
```

- Leverages informal data
- Cuts down the search space

Challenge 1 : different levels of abstraction

- **Solution:** translate into *proof sketches*

We know that $\text{gcd}(a, b) \cdot \text{lcm}(a, b) = ab$,
hence $10 \cdot 280 = n \cdot 40$.

Then $n = 10 \cdot 280 / 40 = 70$

completing the proof. ■

Informal Proof



```
have c1: "10*280 = n*40"  
using assms
```

```
<proof>
```

```
then have c2: "n = 10*280/40"
```

```
<proof>
```

```
then show ?thesis
```

```
<proof>
```

Formal Proof *Sketch*

Challenge 1 : different levels of abstraction

- **Solution:** translate into *proof sketches*
- In-line comments show the alignment between the informal and formal proof.

Informal Statement: Show that for any real number a , $10a \leq 28a^2 + 1$.

Informal Proof:

It suffices to show $0 \leq 28a^2 - 10a + 1$. First, consider completing the square for $28a^2 - 10a$ and observe that $(a - \frac{5}{28})^2 = a^2 - \frac{10}{28}a + (5/28)^2$. Since $0 \leq (a - \frac{5}{28})^2$, we get $0 \leq a^2 - \frac{10}{28}a + (5/28)^2$. Multiplying by 28 and simplifying gives $0 \leq 28a^2 - 10a + (25/28)$. Since $25/28 < 1$, the result follows.

Formal Statement:

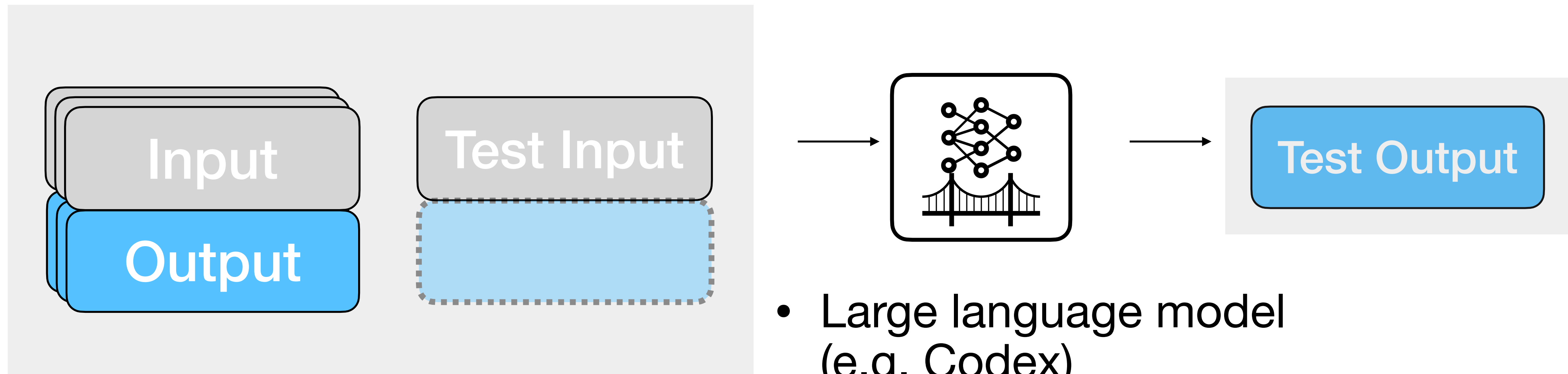
```
theorem algebra_binomnegdiscrineq_10alt28asqp1:  
  fixes a :: real  
  shows "10 * a ≤ 28 * a^2 + 1"
```

Formal Proof Sketch:

```
proof - (* it suffices to show 0 ≤ 28a^2 - 10a + 1 *)  
  have c0: "0 ≤ 28a^2 - 10a + 1"  
  proof - (* observe that (a - (5/28))^2 = a^2 - (10/28)a + (5/28)^2 *)  
    have c1: "(a - (5/28))^2 = a^2 - 10/28a + (5/28)^2" <...>  
    (* we get 0 ≤ a^2 - (10/28)a + (5/28)^2 *)  
    have c2: "0 ≤ a^2 - 10/28a + (5/28)^2" using c1 <...>  
    (* Multiplying by 28 and simplifying gives 0 ≤ 28a^2 - 10a + (25/28) *)  
    have c3: "0 ≤ 28a^2 - 10a + 28((5/28)^2)" using c2 <...>  
    have c4: "0 ≤ 28a^2 - 10a + 28((5/28)*(5/28))" using c3 <...>  
    have c5: "0 ≤ 28a^2 - 10a + (25/28)" using c4 <...>  
    (* Since 25/28 < 1, the result follows. *)  
    show ?thesis using c5 <...>  
  qed  
  show ?thesis <...>  
qed
```


Challenge 2 : no parallel data

- **Solution:** few-shot in-context learning



Challenge 2 : few-shot sketching

Informal Statement: Show that for any real number a , $10a \leq 28a^2 + 1$.

Informal Proof:

It suffices to show $0 \leq 28a^2 - 10a + 1$. First, consider completing the square for $28a^2 - 10a$ and observe that $(a - \frac{5}{28})^2 = a^2 - \frac{10}{28}a + (5/28)^2$. Since $0 \leq (a - \frac{5}{28})^2$, we get $0 \leq a^2 - \frac{10}{28}a + (5/28)^2$. Multiplying by 28 and simplifying gives $0 \leq 28a^2 - 10a + (25/28)$. Since $25/28 < 1$, the result follows.

Formal Proof Sketch:

theorem algebra_binomnegdiscrineq_10alt28asqp1:

fixes a :: real

shows "10 * a ≤ 28 * a^2 + 1"

proof - (* it suffices to show $0 \leq 28a^2 - 10a + 1$ *)

have c0: "0 ≤ 28a^2 - 10a + 1"

proof - (* observe that $(a - (5/28))^2 = a^2 - (10/28)a + (5/28)^2$ *)

have c1: "(a - (5/28))^2 = a^2 - 10/28a + (5/28)^2" <...>

(* we get $0 \leq a^2 - (10/28)a + (5/28)^2$ *)

have c2: "0 ≤ a^2 - 10/28a + (5/28)^2" **using** c1 <...>

(* Multiplying by 28 and simplifying gives $0 \leq 28a^2 - 10a + (25/28)$ *)

have c3: "0 ≤ 28a^2 - 10a + 28((5/28)^2)" **using** c2 <...>

have c4: "0 ≤ 28a^2 - 10a + 28((5/28)*(5/28))" **using** c3 <...>

have c5: "0 ≤ 28a^2 - 10a + (25/28)" **using** c4 <...>

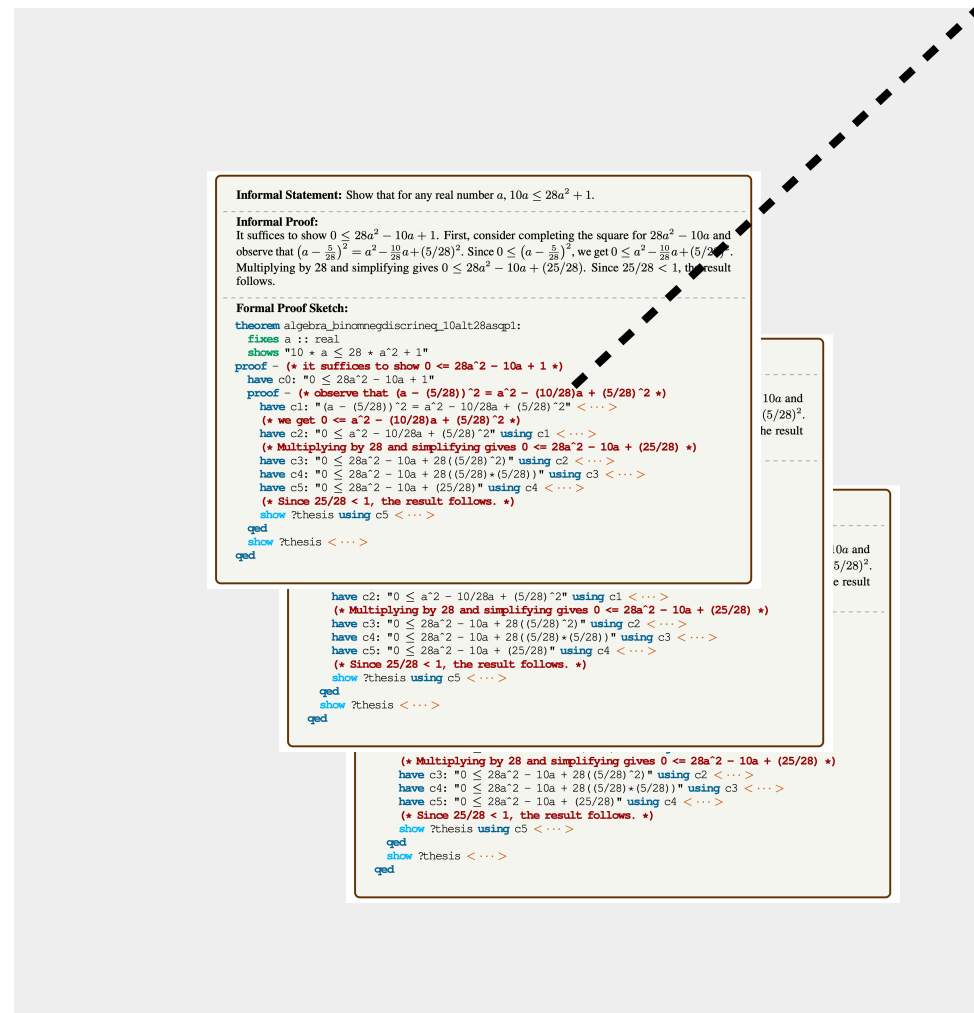
(* Since $25/28 < 1$, the result follows. *)

show ?thesis **using** c5 <...>

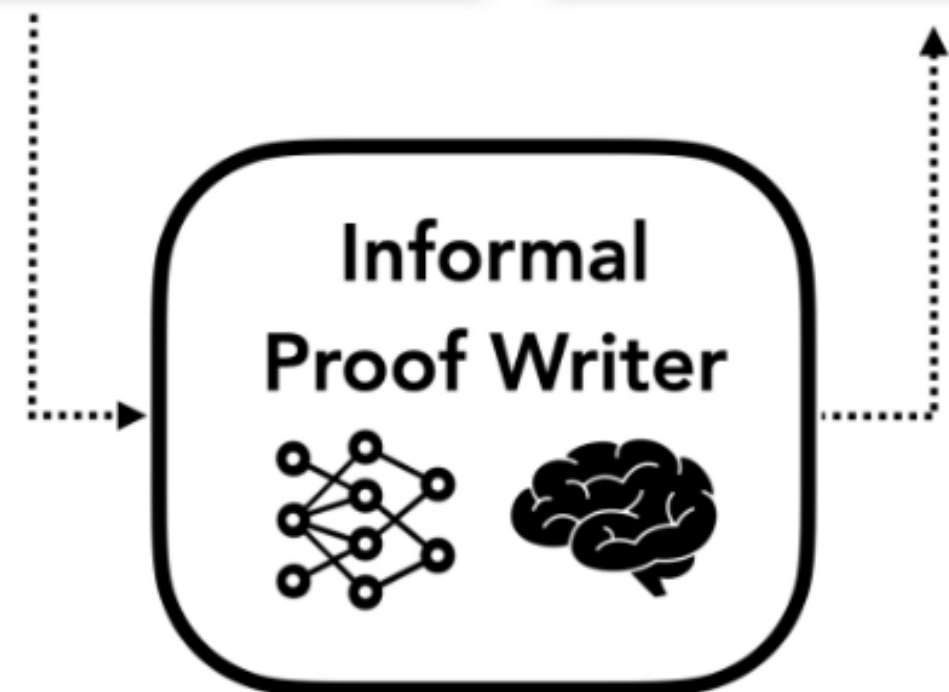
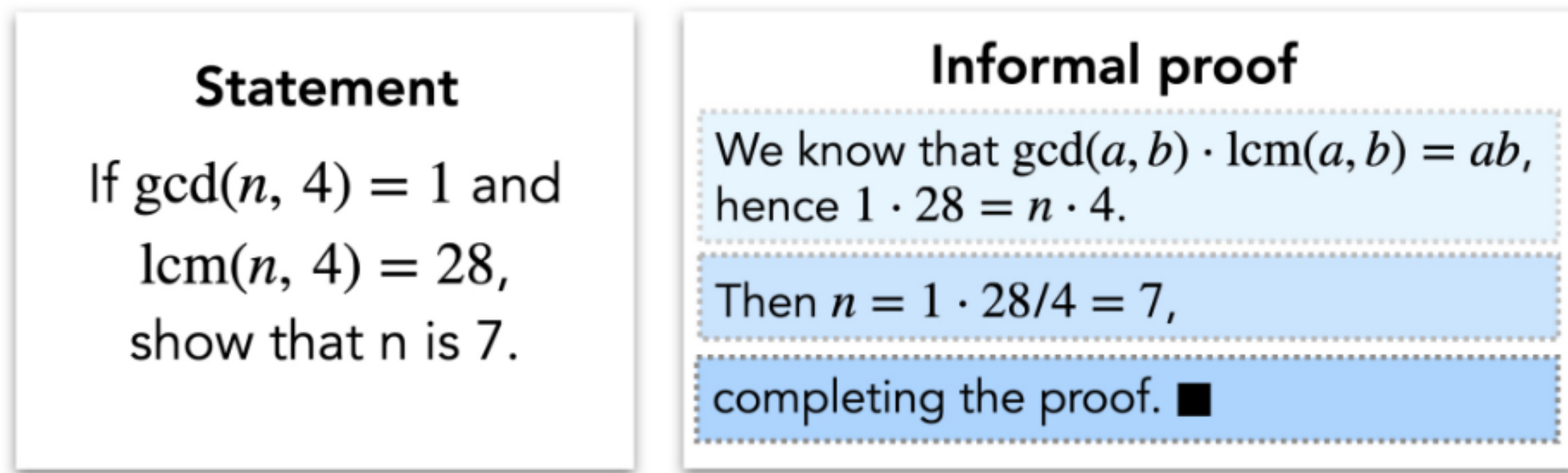
qed

show ?thesis <...>

qed



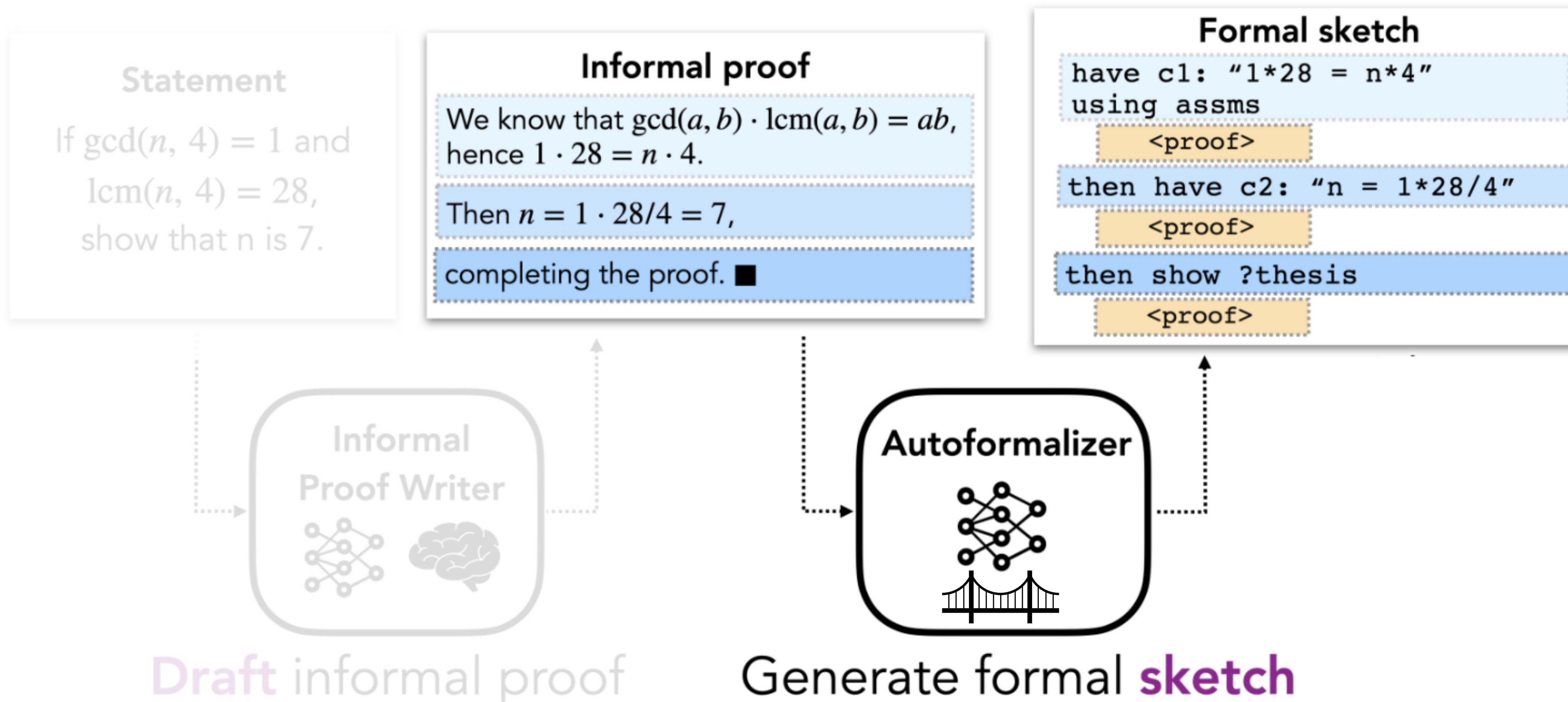
Draft, sketch, prove



Draft informal proof

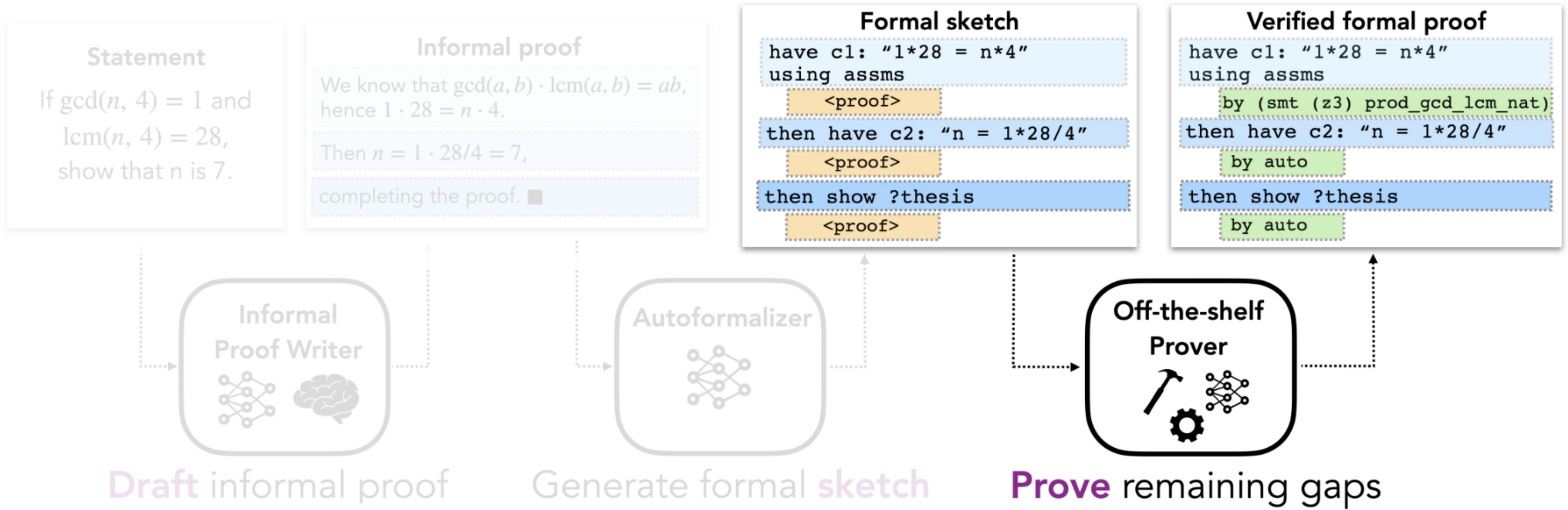
- Human-written proof
- Large language model (Minerva)

Draft, **sketch**, prove



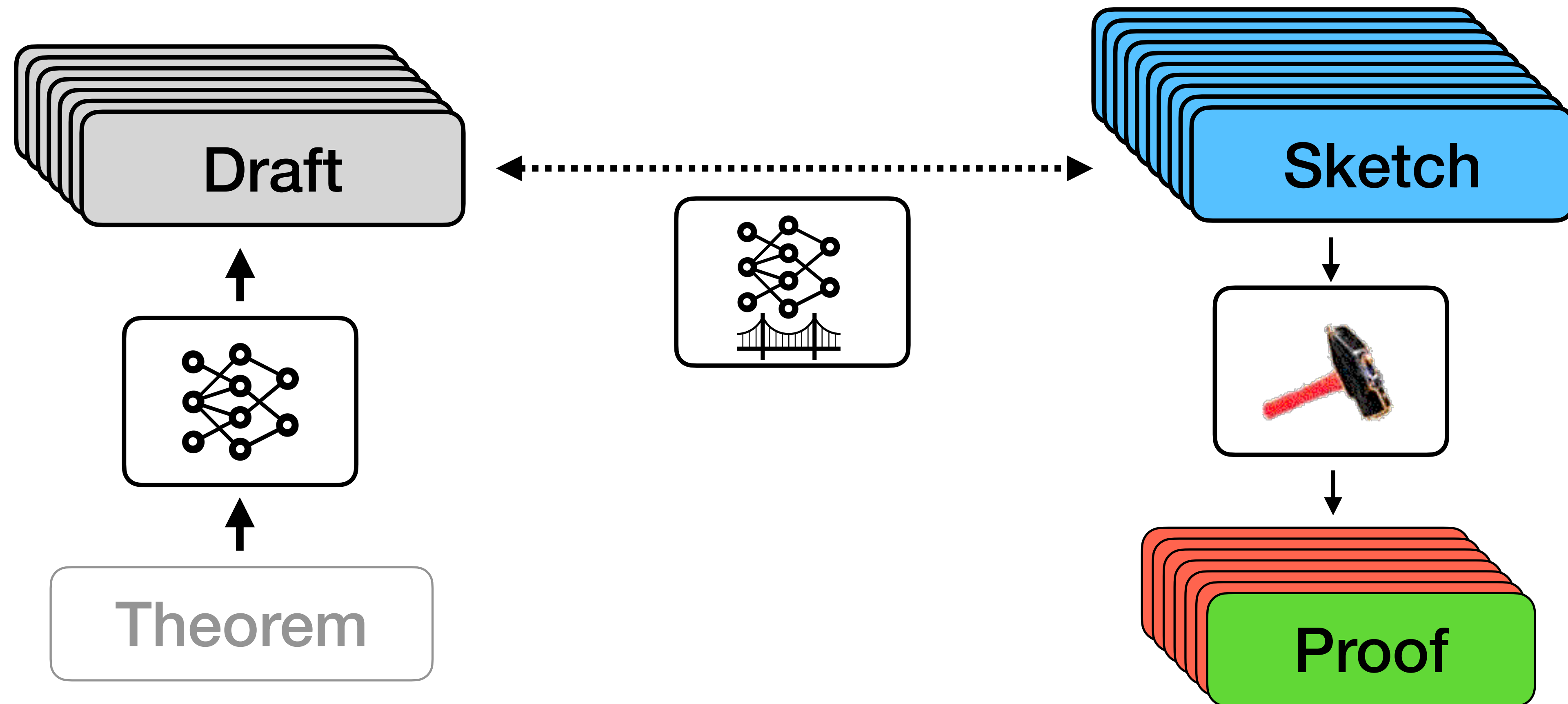
- Large language model (Codex)

Draft, sketch, **prove**



- Sledgehammer + heuristics

Proof search with draft, sketch, prove



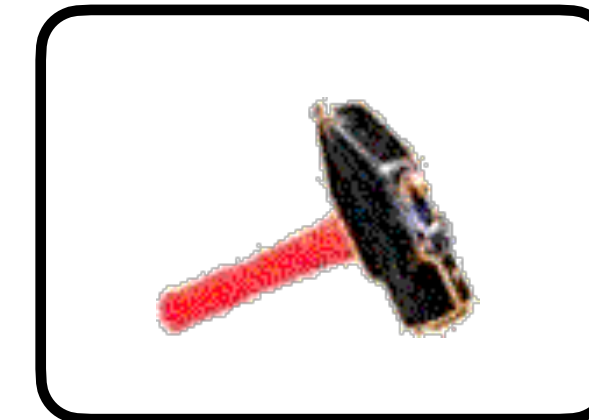
Experiments | miniF2F [Zheng et al 2022]

- 488 high-school competition problems (AMC, AIME, IMO, ...)



- 244 validation, 244 test

- Isabelle proof assistant



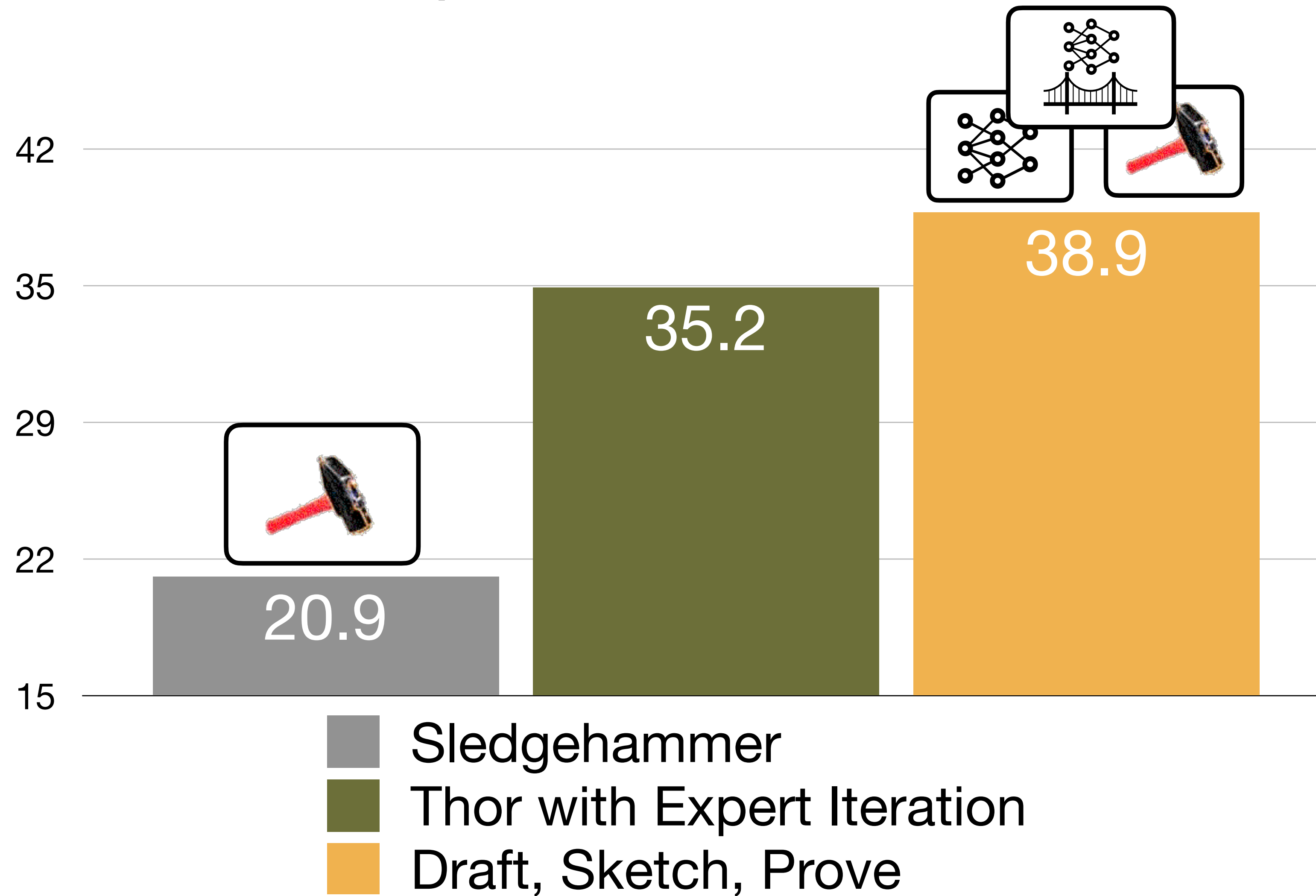
- Baselines:

- Sledgehammer
- THOR with Expert Iteration [Wu et al 2022]

Experiments | miniF2F

- Human informal draft: 1 human draft x 100 sketches/draft
- LLM informal drafts: 100 drafts x 1 sketch / draft
- Prompting:
 - Write 20 sketch examples from miniF2F
 - Sample 3 examples and put in prompt

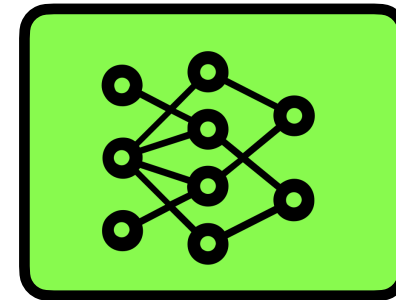
Top-level results



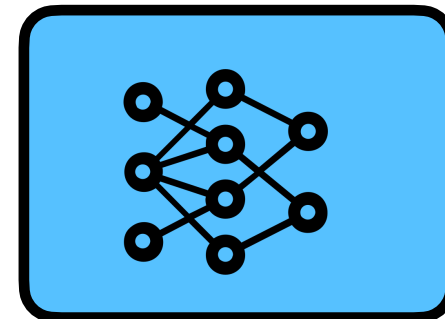
International Math Olympiad



Draft
(Minerva)



Sketch
(Codex)



Formal Prover
(Sledgehammer)



Statement: Prove that the fraction $\frac{21n+4}{14n+3}$ is irreducible for every natural number n .

Informal Proof (Minerva 62B):

We must show that $\gcd(21n + 4, 14n + 3) = 1$.

The Euclidean algorithm gives

$$21n + 4 = 1 \cdot (14n + 3) + 7n + 1$$

$$14n + 3 = 2 \cdot (7n + 1) + 1.$$

Since $\gcd(7n + 1, 1) = 1$, we have $\gcd(21n + 4, 14n + 3) = 1$.

Formal Proof:

```
theorem imo_1959_p1:
```

```
  fixes n :: nat
```

```
  shows "gcd (21*n + 4) (14*n + 3) = 1"
```

```
proof -
```

```
(* The Euclidean algorithm gives
```

```
21n+4=1\cdot(14n+3)+7n+1
```

```
14n+3=2\cdot(7n+1)+1. *)
```

```
have c0: "21*n + 4 = 1*(14*n + 3) + 7*n + 1"
```

```
  by auto [ATP]
```

```
have c1: "14*n + 3 = 2*(7*n + 1) + 1" using c0
```

```
  by auto [ATP]
```

```
(* Since \gcd(7n+1,1)=1, we have \gcd(21n+4,14n+3)=1. *)
```

```
then have "gcd (7*n + 1) 1 = 1"
```

```
  using c1
```

```
  by auto [ATP]
```

```
then have "gcd (21*n + 4) (14*n + 3) = 1"
```

```
  using c1
```

```
  by (smt (z3) BitM_plus_one ab_semigroup_add_class.add_ac(1)
```

```
  add.assoc c0 gcd.commute gcd_add2 gcd_add_mult mult_numeral_1
```

```
  numeral_One numeral_eq_Suc numerals(1) semiring_norm(3)) [ATP]
```

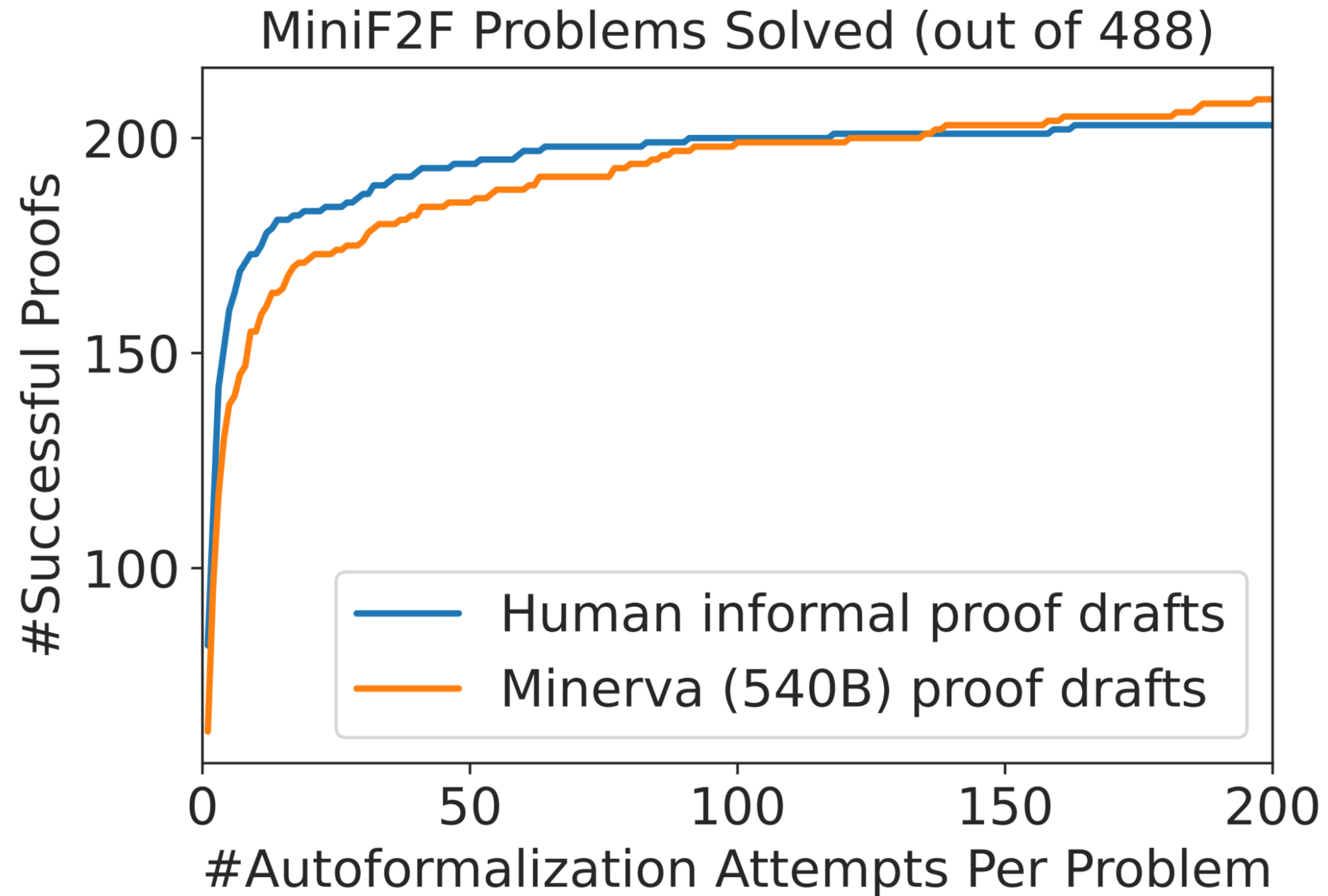
```
then show ?thesis
```

```
  using c1
```

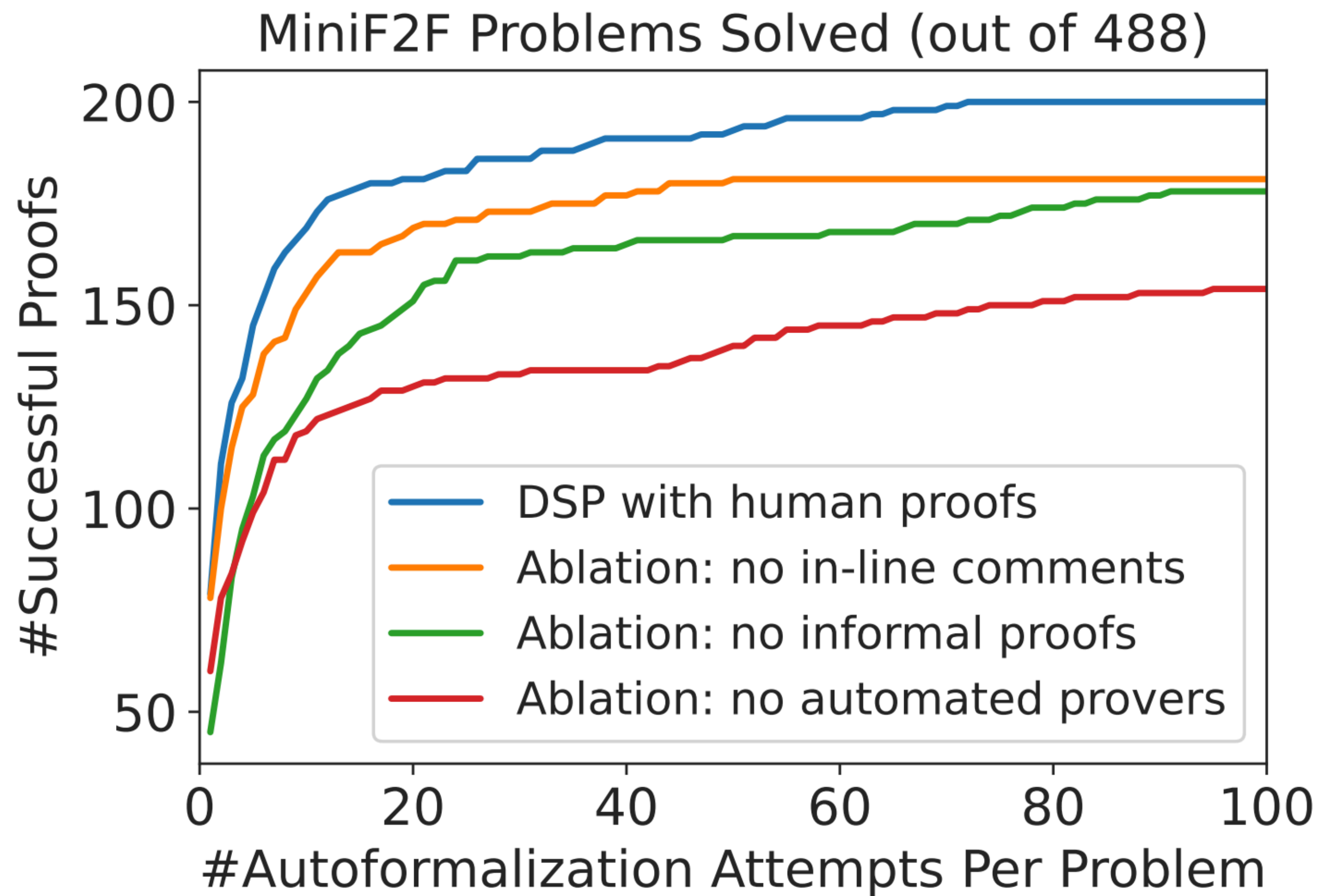
```
  by blast [ATP]
```

```
qed
```

Increasing search budget



Ablations



Model vs. human informal draft

Statement: For any integer a , show that $a^2 \equiv 0 \pmod{4}$ or $a^2 \equiv 1 \pmod{4}$.

Informal Proof (Minerva 62B):

Let $a = 2k$ for some integer k .

Then $a^2 = 4k^2$.

Since $4k^2$ is divisible by 4, we have $a^2 \equiv 0 \pmod{4}$.

Now, let $a = 2k + 1$ for some integer k .

Then $a^2 = 4k^2 + 4k + 1 = 4(k^2 + k) + 1$.

Since $k^2 + k$ is an integer, $4(k^2 + k) + 1$ is not divisible by 4.

Thus, $a^2 \equiv 1 \pmod{4}$.

Statement: For any integer a , show that $a^2 \equiv 0 \pmod{4}$ or $a^2 \equiv 1 \pmod{4}$.

Informal Proof (Human):

$a \pmod{4} \in \{0, 1, 2, 3\}$.

Using that for any natural number k , $a \equiv b \pmod{4}$ implies $a^k \equiv b^k \pmod{4}$, we have $a^2 \pmod{4} \in \{0, 1, 4, 9\}$.

Since $4 \equiv 0 \pmod{4}$ and $9 \equiv 1 \pmod{4}$, the result follows.

Model vs. human informal draft

Formal Proof:

```
theorem numbertheory_sqmod4in01d:
  fixes a :: int
  shows "(a^2 mod 4 = 0)
  \<or> (a^2 mod 4 = 1)"
proof (cases "even a")
case True
  (* Let a=2k for some integer k.
  Then a^2=4k^2.
  Since 4k^2 is divisible by 4,
  we have a^2 \equiv 0 \pmod{4}.* *)
  then obtain k where "a=2*k"
  using evenE
  <ATP> by auto </ATP>
  then have "a^2 = 4*k^2"
  unfolding power2_eq_square
  <ATP> by auto </ATP>
  then have "a^2 mod 4 = 0"
  <ATP> by auto </ATP>
  then show ?thesis
  <ATP> by auto </ATP>
next
case False
  (* Now, let a=2k+1 for some integer k.
  Then a^2=4k^2+4k+1=4(k^2+k)+1.
  Since k^2+k is an integer,
  4(k^2+k)+1 is not divisible by 4.
  Thus, a^2 \equiv 1 \pmod{4}.* *)
  then obtain k where "a=2*k+1"
  using oddE <ATP> by auto </ATP>
  then have "a^2 = 4*k^2+4*k+1"
  unfolding power2_eq_square <ATP>
  by (auto simp: field_simps) </ATP>
  then have "a^2 mod 4 = 1"
  <ATP> by presburger </ATP>
  then show ?thesis
  <ATP> by auto </ATP>
qed
```

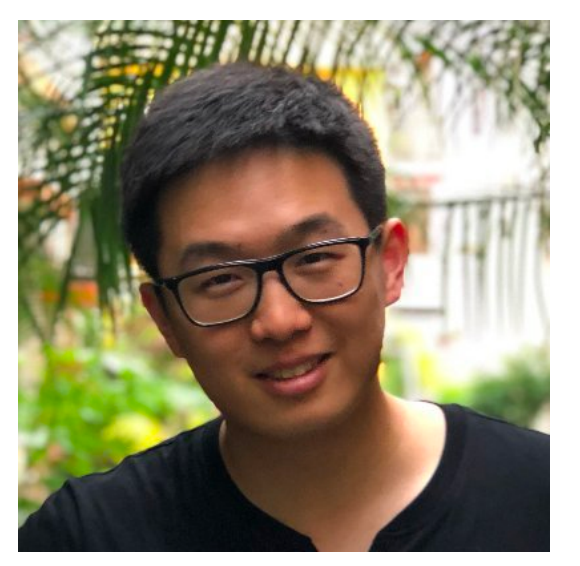
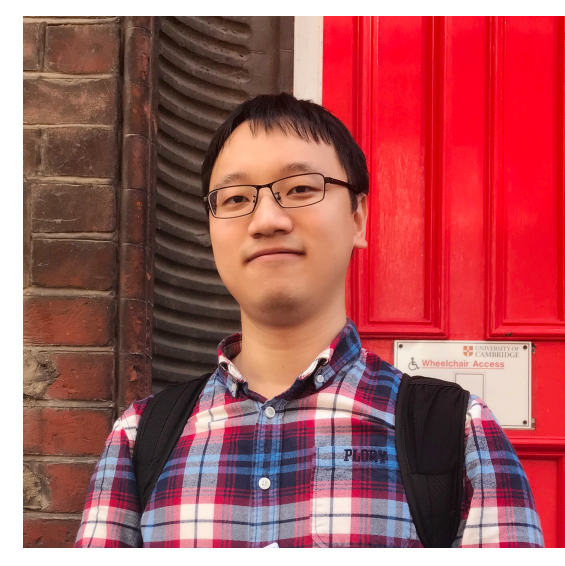
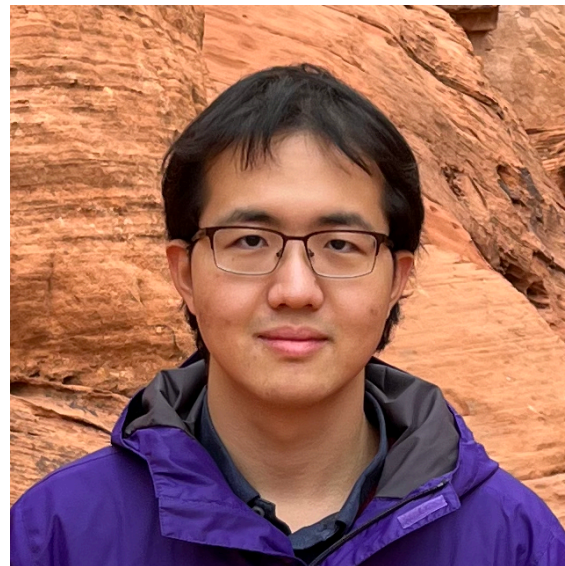
Formal Proof:

```
theorem numbertheory_sqmod4in01d:
  fixes a :: int
  shows "(a^2 mod 4 = 0)
  \<or> (a^2 mod 4 = 1)"
proof -
  (* a \pmod 4 \in {0, 1, 2, 3}. *)
  have c0: "a mod 4 \<in> {0, 1, 2, 3}"
  <ATP> by auto </ATP>
  (* Using that for any natural number k,
  a \equiv b \pmod 4 implies
  a^k \equiv b^k \pmod 4, *)
  have "a^2 mod 4 = (a mod 4)\<sup>
  2 mod 4" <ATP> by (smt (z3)
  numeral_eq_Suc power_mod) </ATP>
  (* we have
  a^2 \pmod 4 \in {0, 1, 4, 9}. *)
  also have "... \<in> {0, 1, 4, 9}"
  using c0
  <ATP> by auto </ATP>
  (* Since 4 \equiv 0 \pmod 4 and
  9 \equiv 1 \pmod 4,
  the result follows. *)
  finally show ?thesis
  <ATP> by auto </ATP>
qed
```

Conclusion

- Draft informal proofs and sketch formal proofs
 - Combines flexibility of language models with formal systems
- Opens many possibilities and questions
 - Memorization
 - Alternative ways of integrating language models
 - Exploration and reinforcement learning

Thank you



Appendix: full results

Success rate	miniF2F-valid	miniF2F-test
<i>Baselines</i>		
Sledgehammer	9.9%	10.4%
Sledgehammer + heuristics	18.0%	20.9%
Thor (Jiang et al., 2022)	28.3%	29.9%
Thor + expert iteration (Wu et al., 2022)	37.3%	35.2%
<i>Draft, Sketch, and Prove</i>		
Human informal proof	42.6%	39.3%
Codex informal proof	40.6%	35.3%
8B Minerva informal proof	40.6%	35.3%
62B Minerva informal proof	43.9%	37.7%
540B Minerva informal proof	42.6%	38.9%
<i>Ablations (with human informal statements and proofs)</i>		
– In-line comments	37.7% (−4.9%)	36.5% (−2.8%)
– Informal proofs	38.9% (−3.7%)	34.0% (−5.3%)
– Automated provers	32.8% (−9.8%)	30.3% (−9.0%)